CS401 - Problem Set 8

- 1. (a) Without oracles, we proved $NP \subseteq PSPACE$ using simulation. This is because, in our proof, the PSPACE machine simulates what the NP machine does. Show that the proof still works for oracular machines. That is, prove that for every language O, $NP^{O} \subseteq PSPACE^{O}$.
 - (b) The example from part (a) should help you see why simulation proofs "relativize." Please explain more generally why, if you prove $\mathbf{A} \subseteq \mathbf{B}$ using a simulation, then for any oracle $\mathbf{O}, \mathbf{A}^{\mathbf{O}} \subseteq \mathbf{B}^{\mathbf{O}}$? (This need not be a formal proof try to make your explanation as simple as possible.)
- 2. (a) Prove (or sketch a proof) that for every language O, $\mathsf{DTIME}(n)^{\mathsf{O}} \subseteq \mathsf{DTIME}(n^{1.5})^{\mathsf{O}}$.
 - (b) The example from part (a) should help you see why diagonalization proofs "relativize." Please explain more generally why, if you prove $\mathbf{A} \subseteq \mathbf{B}$ using a diagonalization, then for any oracle $\mathbf{O}, \mathbf{A}^{\mathbf{O}} \subseteq \mathbf{B}^{\mathbf{O}}$? (This need not be a formal proof try to make your explanation as simple as possible. Think about what properties of TMs were needed for diagonalization, and argue those properties are still present for oracular TMs.)
- 3. [Problem from Class] Prove $coNP \subseteq P^{SAT}$, where recall SAT is the language of Boolean formulas that have a satisfying assignment.