CS401 - Problem Set 2

- 1. If we discussed revisions to your learning plan, please make those changes.
- 2. In this problem, you'll explore how the base used in writing down the language affects the complexity class the language is in.
 - (a) For $S \subseteq \mathbb{N}$, let $L_b^S = \{n_b : n \in S\}$, where $n_b \in \{0, 1, \dots, b-1\}^*$ is a string that represents the natural number n written in base b. Prove that for b > 2, if $L_b^S \in \mathbf{P}$, then $L_2^S \in \mathbf{P}$.
 - (b) Explain (using as plain, accessible language as possible) why your proof in part b justifies our choice to only consider binary languages in this class (at least in cases where we have at least polynomial time as a resource).
 - (c) * Prove that the following language is in \mathbf{P} :

UNARYFACTORING = {
$$\langle n_{\text{unary}}, l_{\text{unary}}, k_{\text{unary}} \rangle$$
 : $\exists j \in \mathbb{N}$ s.t. $l \leq j \leq k$ and j divides n }. (1)

In this case, n_{unary} means the number *n* represented in unary. A unary string is an element of $\{1\}^*$ where a natural number is represented using a string containing that number of ones. So for example, 2 is represented in unary as 11, and 5 is represented in unary as 11111. Then $\langle 1111111111, 1111, 111111\rangle \in \mathsf{UNARYFACTORING}$ because in base 10, this corresponds to the sequence < 10, 4, 7 >. Since 5 divides 10 and $4 \le 5 \le 7$, this tuple satisfies the conditions, and so is in the language. On the other hand, $\langle 1111111111, 111111\rangle \notin \mathsf{UNARYFACTORING}$ because there is no number between 4 and 7 that divides 9.

- (d) It is unknown whether the problem of factoring in base 2 is in \mathbf{P} . However, in part (c) you prove that UNARYFACTORING $\in \mathbf{P}$. Why can't we use the same argument as in part (a) to convert base-2 factoring into UNARYFACTORING, and thus prove factoring is in \mathbf{P} ?
- 3. (Extra Practice Problem) Let $L_{\triangle} = \{G : G \text{ contains a triangle}\}$. Prove L_{\triangle} is in **P**.
- 4. (Moved to Pset 3)
 - (a) A unary language L is a subset of $\{1\}^*$. Let NP_U be the set of unary languages that are also in NP . Prove that if $\mathsf{NP}_U \subseteq \mathsf{P}$, then $\mathsf{EXP} = \mathsf{NEXP}$, where NEXP is defined similarly to NP except now the TM can run for exponential time in the size of the input, and the witness u can be of exponential size in the size of the input.
 - (b) EXP vs NEXP is the exponential time version of the P vs NP question. Based on the result in part (a), how does the EXP vs NEXP question relate to the P vs NP question? Is this relationship surprising?

- 5. (You almost, but not quite, have the tools to solve this. I'll move it to next pset, but I encourage you to think about it or try to solve it on this pset.) Prove if $L' \in NP$ and $L \leq_p L'$, then $L \in NP$.
- 6. Let $L_{ksort} = \{ \langle x, k \rangle : x \text{ is a list of unsorted integers, and the } k^{\text{th}} \text{ largest integer is odd} \}$. Prove L_{ksort} is in **P**.