

Divide & Conquer Structure

Merge Sort

Input : Integer array A of length n

Output: Sorted array

```
// Base Case
1 if  $n == 1$  then
2   return  $A$ ;
3 end
// Divide and Conquer
4  $A_1 = \text{MergeSort}(A[1 : n/2]);$ 
5  $A_2 = \text{MergeSort}(A[n/2 + 1 : n]);$ 
// Combine
6  $p_1 = p_2 = 1;$ 
7 for  $i=1$  to  $n$  do
8   if  $A_1[p_1] < A_2[p_2]$  then
9      $A[i] = A_1[p_1];$ 
10     $p_1++;$ 
11  else
12     $A[i] = A_2[p_2];$ 
13     $p_2++;$ 
14  end
15 end
```

Base Case

Preprocessing (none in MergeSort)

Divide input into equal sized pieces +
Recursively solve each piece

Combine solutions
from each
piece to get
ultimate solution

MergeSort

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Output: Sorted array

// Base Case

1 if $n == 1$ then

2 | return A ;

3 end

// Divide and Conquer

4 $A_1 = \text{MergeSort}(A[1 : n/2])$;

5 $A_2 = \text{MergeSort}(A[n/2 + 1 : n])$;

// Combine

6 $p_1 = p_2 = 1$;

7 for $i=1$ to n do

8 | if $A_1[p_1] < A_2[p_2]$ then

9 | $A[i] = A_1[p_1]$;

10 | p_1++ ;

11 else

12 | $A[i] = A_2[p_2]$;

13 | p_2++ ;

14 end

15 end

How to prove correct?

Strong Induction

See review

→ videos

if size $(n-1)$ then could use induction

How to analyze runtime?

Recurrence Relation

$T(n)$ = runtime on input size n

$$T(n) = \begin{cases} \dots & \text{if } n \leq 1 \text{ (base cases)} \\ T(n/2) & \text{else} \end{cases}$$