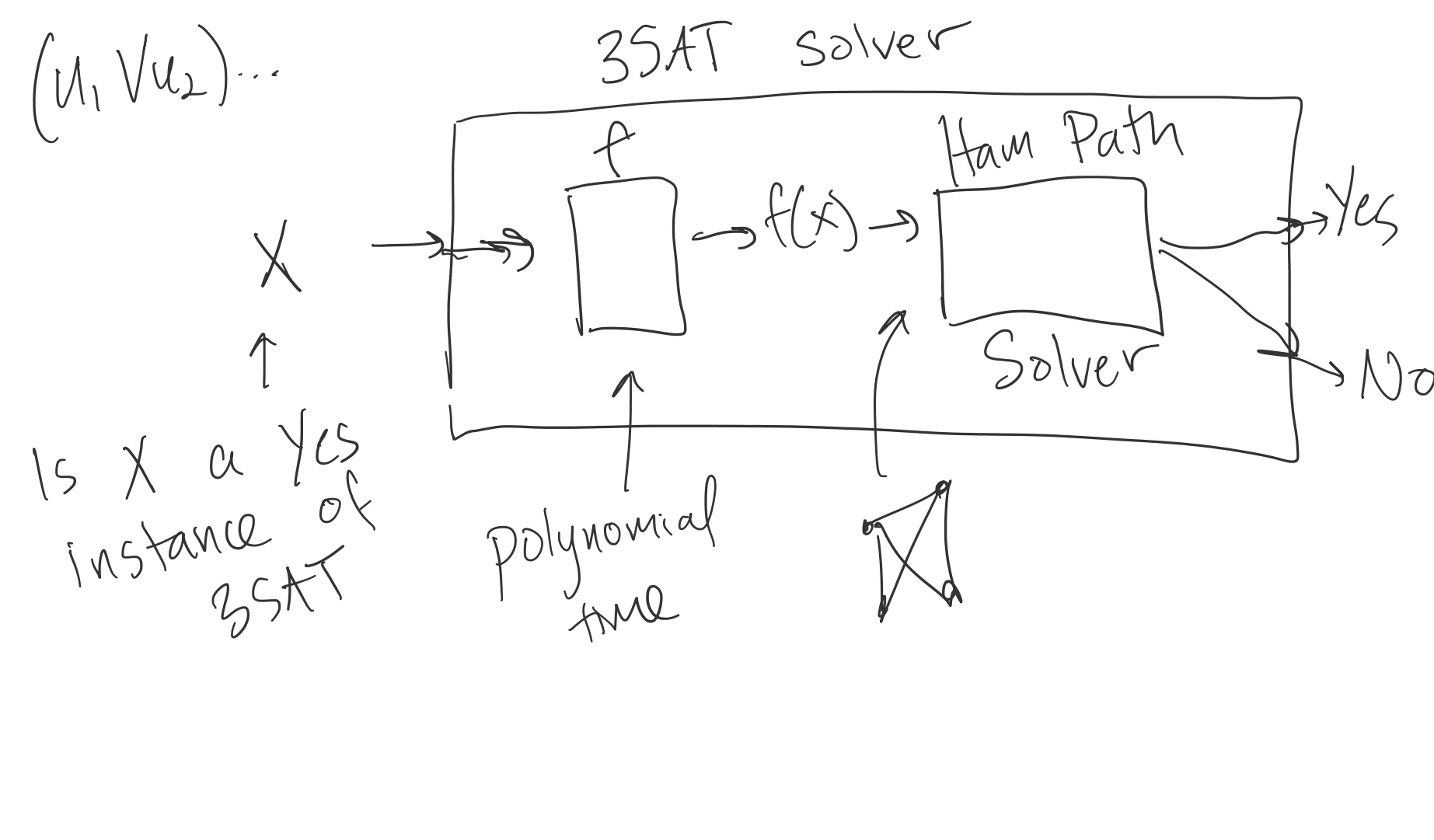
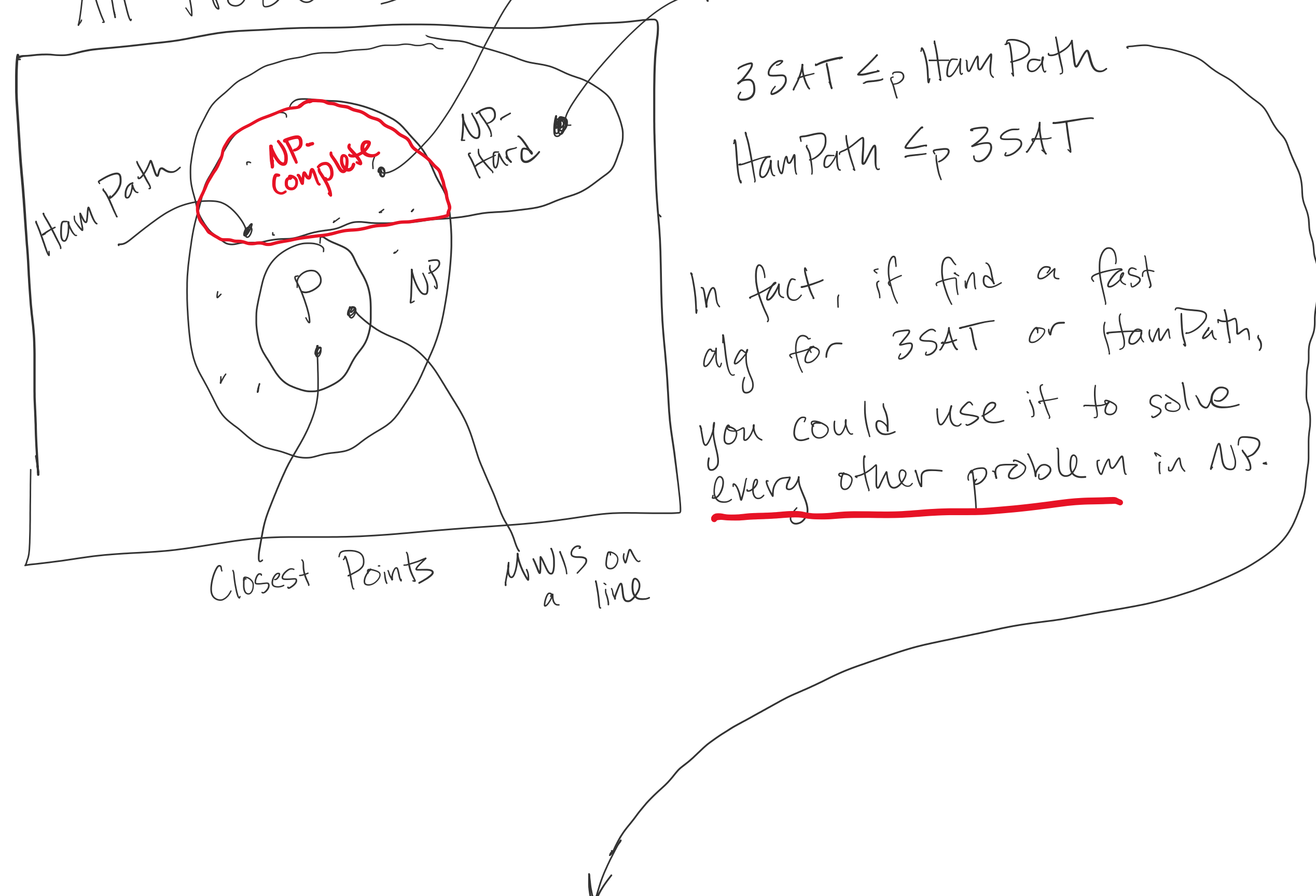


**Goals:**

- Understand meaning of NP-Hard, NP-Complete
- Start describing reduction of 3SAT to HamPath

**Announcements:**

- Go/midhacks2022
- Great exit tickets on Friday- see pinned discussion for my responses
- Don't forget about quiz, midterm revisions! Don't leave for the last minute!



def:

**NP-Hard:**  $Q \in NP\text{-Hard}$  if  $\forall R \in NP, R \leq_p Q$

For all  $R \in NP$



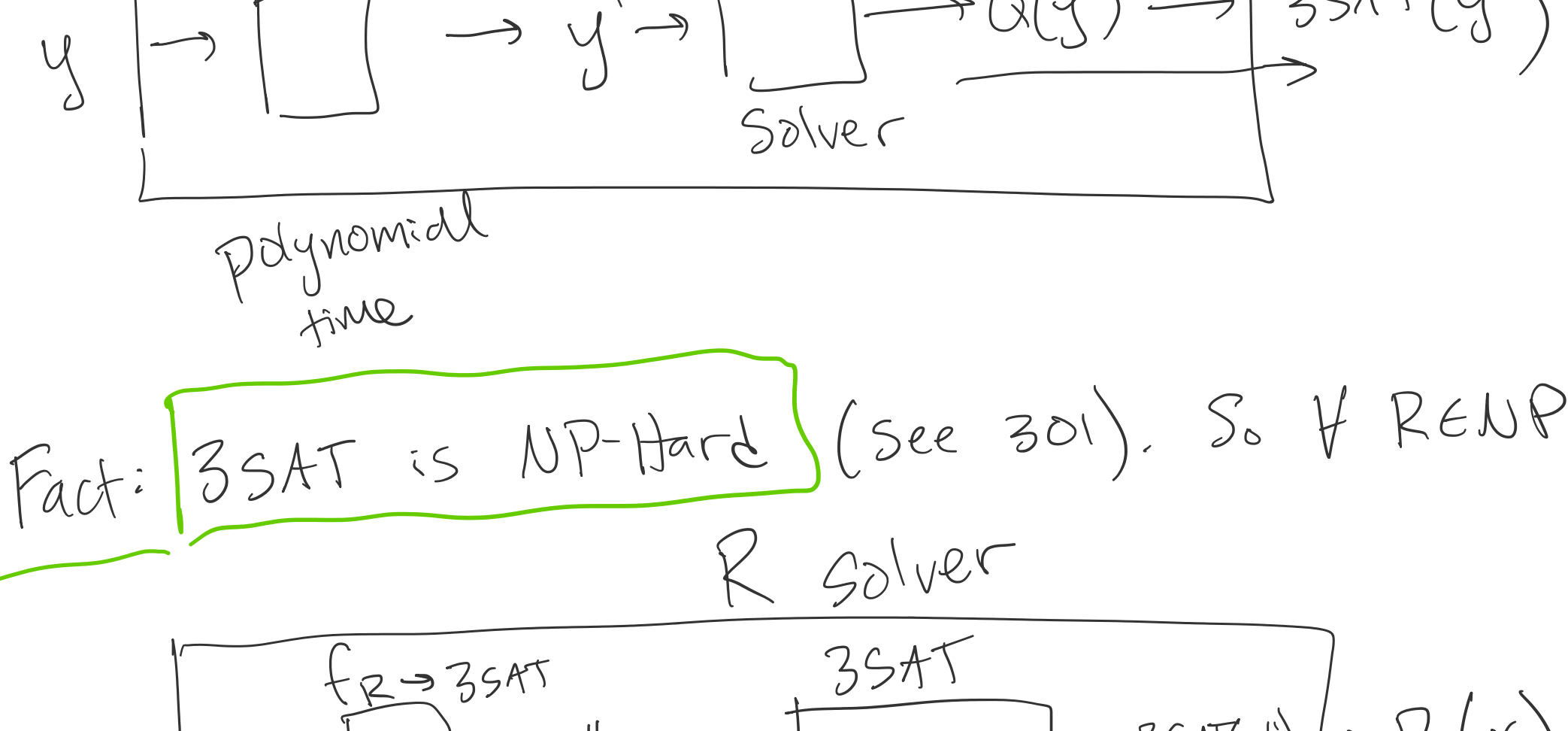
def:

**NP-Complete:** If  $Q \in NP$  and  $Q \in NP\text{-Hard}$  Then  $Q \in NP\text{-Complete}$  How?

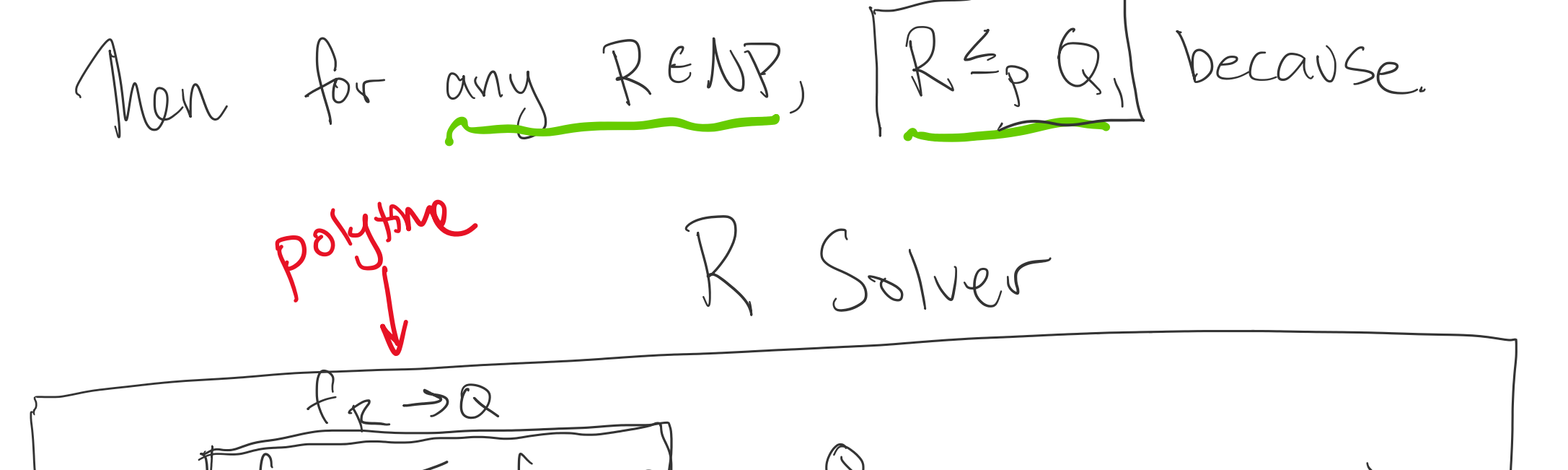
Show  $3\text{-SAT} \leq_p Q$

**Theorem:** If  $3\text{SAT} \leq_p Q$  then  $Q \in NP\text{-Hard}$ :

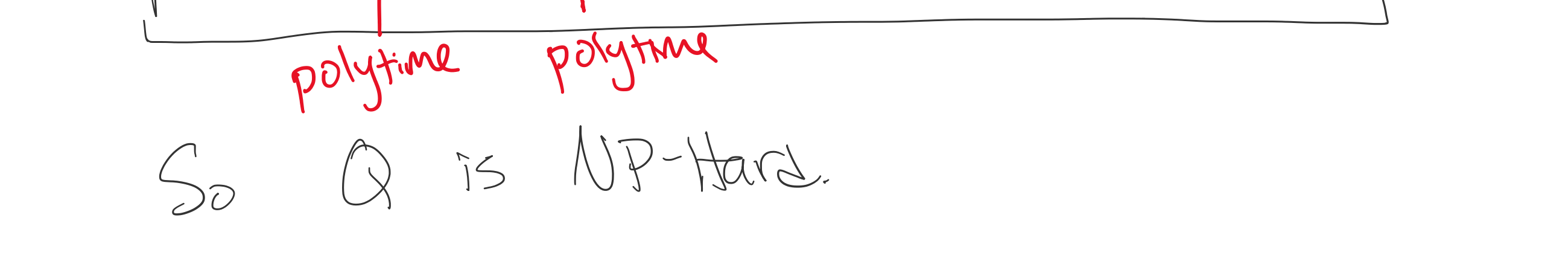
Pf:



**Fact:**  $3\text{SAT}$  is NP-Hard (see 301). So  $\forall R \in NP$



Then for any  $R \in NP, R \leq_p Q$  because



So  $Q$  is NP-Hard.

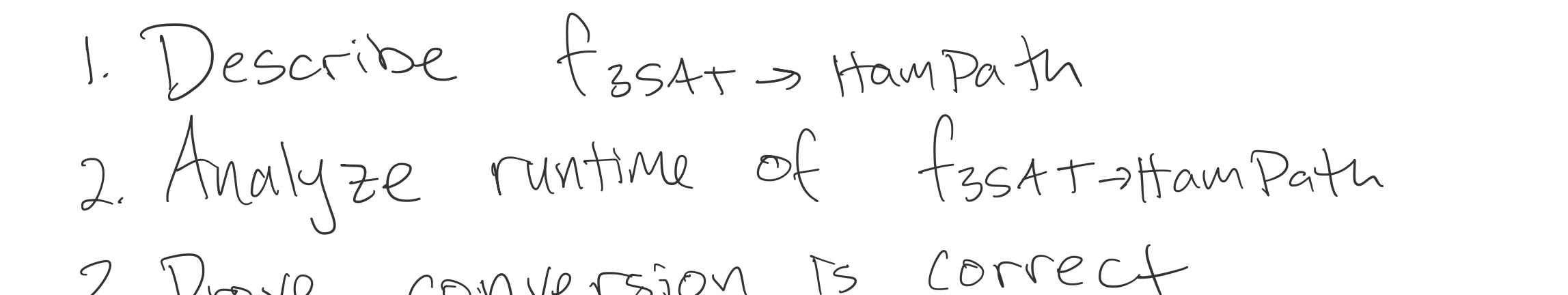
**Group Work**

1. Fill in box to finish proof.
2. Brainstorm ideas for  $3\text{SAT} \leq_p \text{HamPath}$

**Theorem:** Hamiltonian Path is NP-Complete

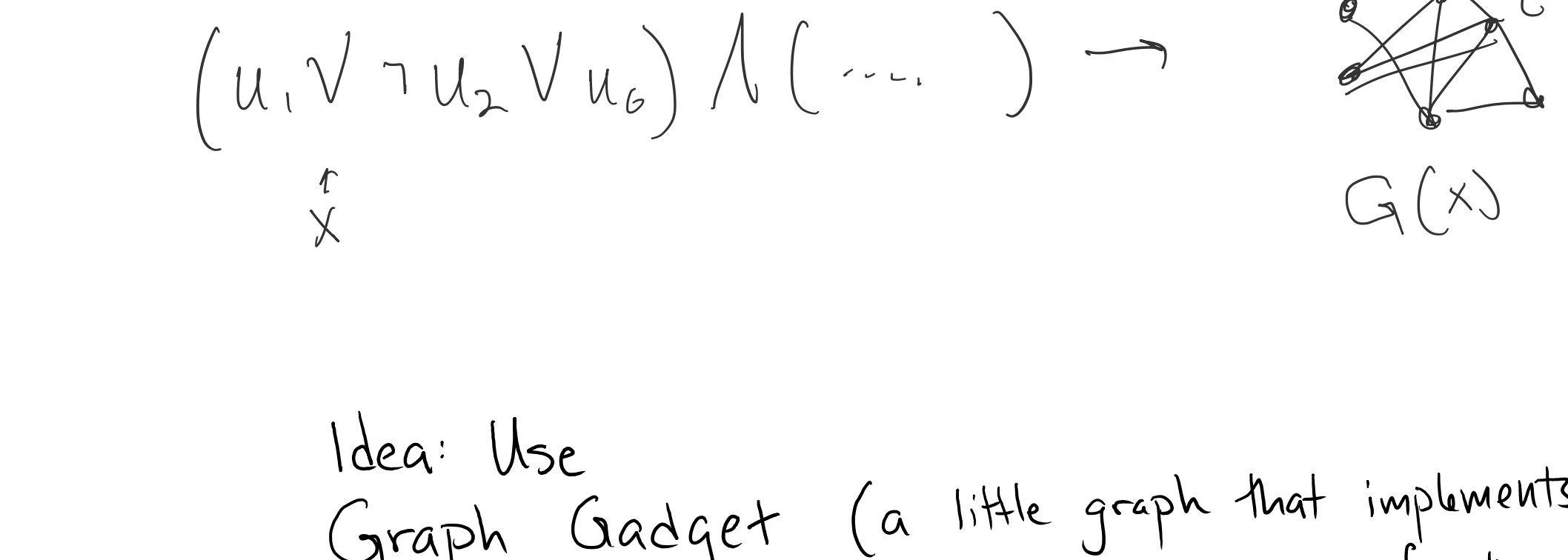
- Pf:
- $\text{HamPath} \in NP$  ✓
  - $\text{HamPath} \in NP\text{-Hard}$

$\hookrightarrow 3\text{SAT} \leq_p \text{HamPath}$

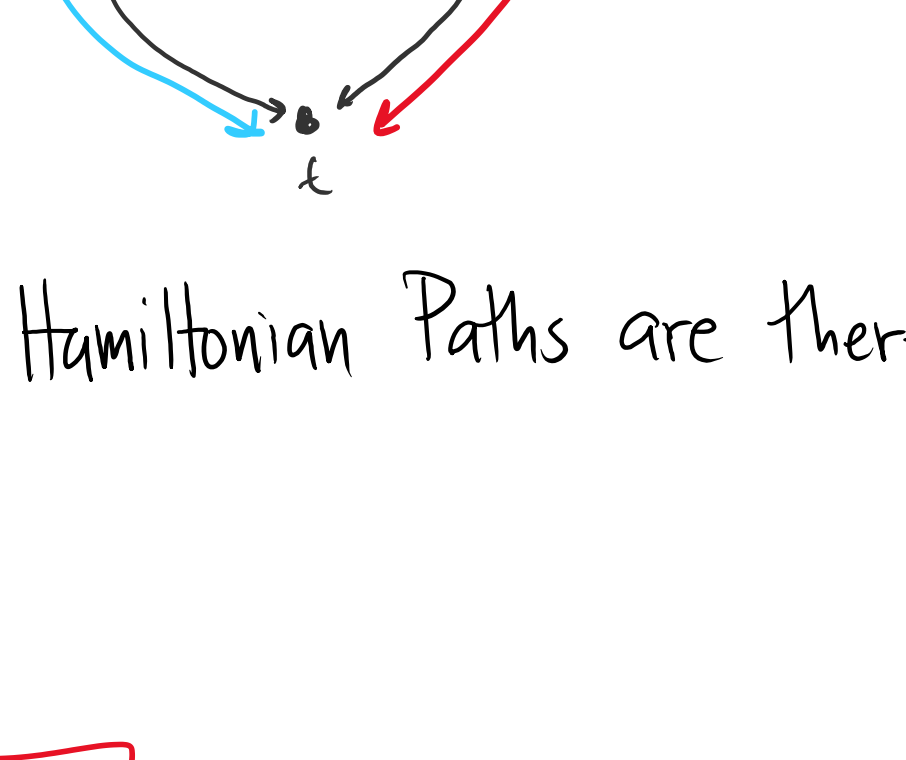


Need to prove this reduction exists

1. Describe  $f_{3\text{SAT} \rightarrow \text{HamPath}}$
2. Analyze runtime of  $f_{3\text{SAT} \rightarrow \text{HamPath}}$
3. Prove conversion is correct
  - a) If  $\text{HamPath}(G(x)) = \text{Yes} \rightarrow 3\text{SAT}(x) = \text{Yes}$
  - b) If  $3\text{SAT}(x) = \text{Yes} \rightarrow \text{HamPath}(G(x)) = \text{Yes}$

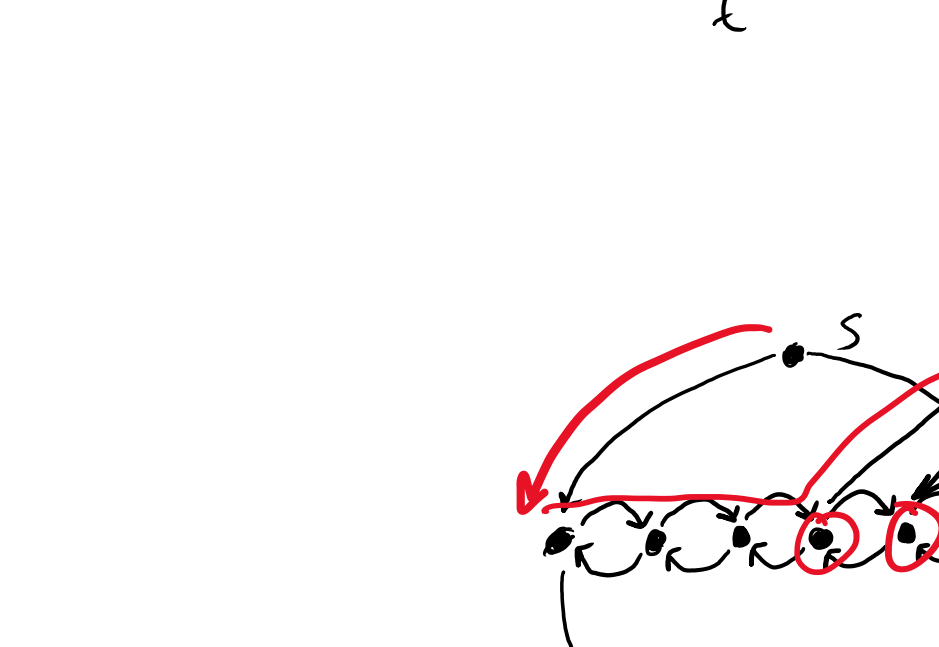


Idea: Use Graph Gadget (a little graph that implements some functionality)



How many Hamiltonian Paths are there from s to t?

- A) 0
- B) 1
- C) 2**
- D) >2



How many Hamiltonian Paths are there from s to t?

- A) 0
- B) 1**
- C) 2
- D) >2

1 path

