Goals:

 Design a dynamic programming algorithm for shortest path

Group work

1.
$$P_{v,i} = \begin{pmatrix} P_{u,i-1} + (u,v) & \text{if shortest path with } i \text{ edges goes} \\ P_{w,i-1} + (w,v) & \text{whith } i \text{ edges goes} \\ P_{w,i-1} + (w,v) & \text{whith } i \text{ edges} \text{$$

$$P_{S,O} = \emptyset$$
 $P_{V,O} = None$
 $V_{V,O} = V_{V,O}$
 $V \in V, \text{ and } V \neq S$

2. Turn into objective function recurrence relation $L(P_{v,i}) = \text{length of shortest path from } s \rightarrow v, \text{ with }$ $L(P_{v,i}) = \text{length of shortest path from } s \rightarrow v, \text{ with }$ $L(P_{v,i}) = \text{min } \sum_{i \in V} \text{min } \sum_{i \in V} \sum_{j \in V} L(P_{u,i-j}) + w(u,v) \sum_{j \in V} L(v,i-j)$ $L(P_{s,o}) = 0 \quad \text{(u,v) of } L(P_{v,o}) = \infty$ $L(P_{s,o}) = 0 \quad \text{(u,v) of } L(P_{v,o}) = \infty$ $L(P_{v,o}) = 0 \quad \text{(u,v) of } L(P_{v,o}) = \infty$

3. Write Pseudocode to Fill in A (don't need to work backward)

Input: Description of an n-vertex graph via an nxn array w, such that w[u,v] contains weight of edge (u,v). (Weight is infinity if no edge and 0 for w[v,v].) Starting vertex s

Output: Array A containing...?

