## CS200 - Problem Set 6

1. Prove Algorithm 1 correctly outputs a list of the prime factors of an integer. The prime factors of n are a list of primes whose product is n. For example for input 60, the algorithm outputs: "2,2,3,5", since 2,3,5 are all prime, and  $2 \times 2 \times 3 \times 5 = 60$ . (To prove this result, you need to prove that the outputs are all prime, and that the product of all of the outputs gives you the input.)

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Input : An integer n such that n \ge 2
  Output: String of the prime factors of n
  /* Recursive Step
                                                               */
1 d = 2;
2 while \neg d | n \operatorname{do}
3 d = d + 1;
4 end
5 if d = n then
     return "n";
6
7 else
     return "d,"+Factor(n/d). // + concatenates strings
8
9 end
                   Algorithm 1: Factor(n)
```

2. Max-Weight-Independent-Set is a problem with many applications, but is used in particular to schedule communications among cell towers. You might study it in CS 302. The following is a part of the correctness proof of the algorithm. Let  $[m] = \{1, 2, 3, \ldots, m\}$ . An independent set S on [m] satisfies:  $S \subseteq [m]$  and if  $i \in S$ , then  $i + 1 \notin S$  and  $i - 1 \notin S$ . Given a function  $f : [n] \to \mathbb{N}$  and  $1 \le m \le n$ , we would like to find the independent set S on [m] such that  $\sum_{i \in S} f(i)$  (the weight) is as large as possible.  $\sum_{i \in S} f(i)$  is the sum of the values of f of all of the elements in the set S. Let  $S_m^*$  be the independent set with maximum weight on [m].

For example, if f(1) = 4, f(2) = 3, f(3) = 2, and f(4) = 5, then  $S_4^* = \{1, 4\}$ , and  $S_3^* = \{1, 3\}$ .

- (a) Prove that if there is a unique max-weight independent set  $S_m^*$  and if  $m \notin S_m^*$ , then  $S_m^* = S_{m-1}^*$ .
- (b) Prove that if there is a unique max-weight independent set  $S_m^*$  and if  $m \in S_m^*$ , then  $S_m^* = \{m\} \cup S_{m-2}^*$ .
- 3. How many bit strings of length 8 contain at least 6 consecutive 0's?
- 4. What is the tightest (smallest) big-O bound that you can prove on  $e^{-x}$ : O(1), O(x), or  $O(2^x)$ . Give a k and C for your response.

- 5. (a) What is the largest number of edges possible in an undirected graph with n vertices and no self-loops. (A self loop is an edge from a vertex back to itself). Please use a counting rule to count possible edges.
  - (b) What is the largest number of edges possible in an undirected graph with n vertices and self-loops.
  - (c) What is the largest number of edges possible in a directed graph (where every edge is of the form (a, b) if a and b are vertices) with n vertices and no self-loops.
  - (d) What is the largest number of edges possible in a directed graph with n vertices and self-loops.
- 6. How long did you spend on this homework?