

FACTORING

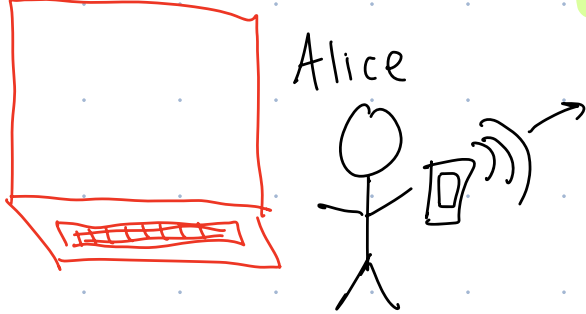
Learning Goals

- Analyze a multi-qubit algorithm
- Become familiar with Shor's factoring algorithm (most famous q. alg)

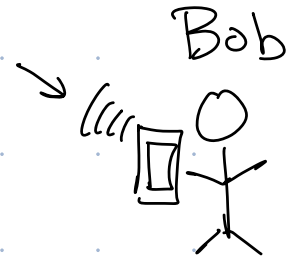
Public Key Crypto (RSA)

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③



①



④



Eve has access to:
If she could:

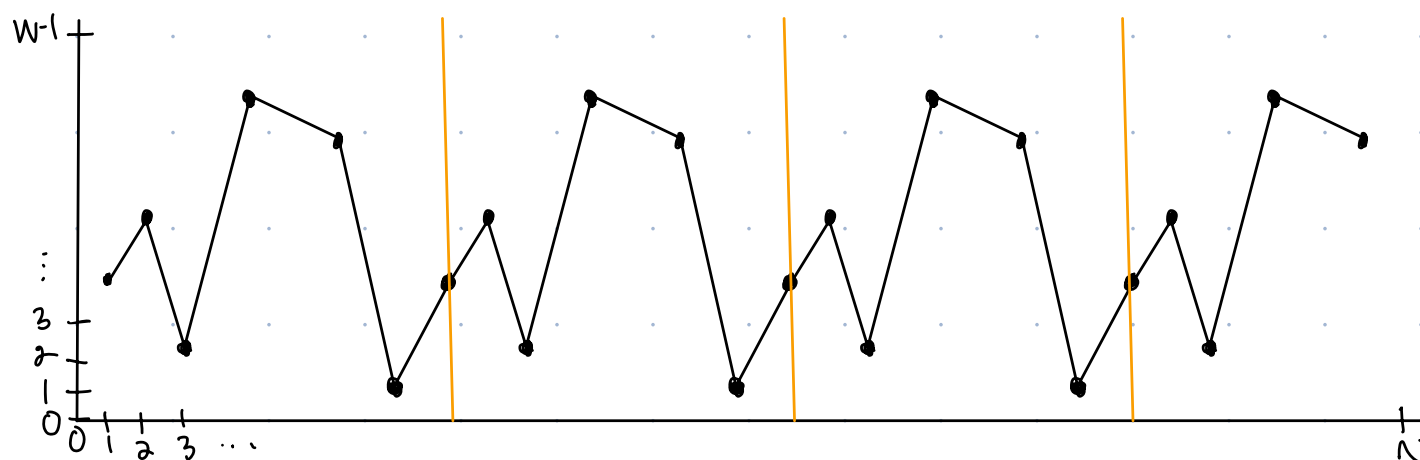
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Factoring reduces (via number theory) to period finding:

Period Finding Problem

Input:

ex:



Output: r

What is the classical query complexity of factoring?

- A) $O(\sqrt{r})$ B) $O(r)$ C) $O(r^2)$ D) $O(n)$

Bits to Digits

Classically, we code using base 10, not binary. We'll do same:

Standard Basis States:

Ambiguity:

If $|3\rangle$ is an N -dimensional state, how many qubits are in the system?

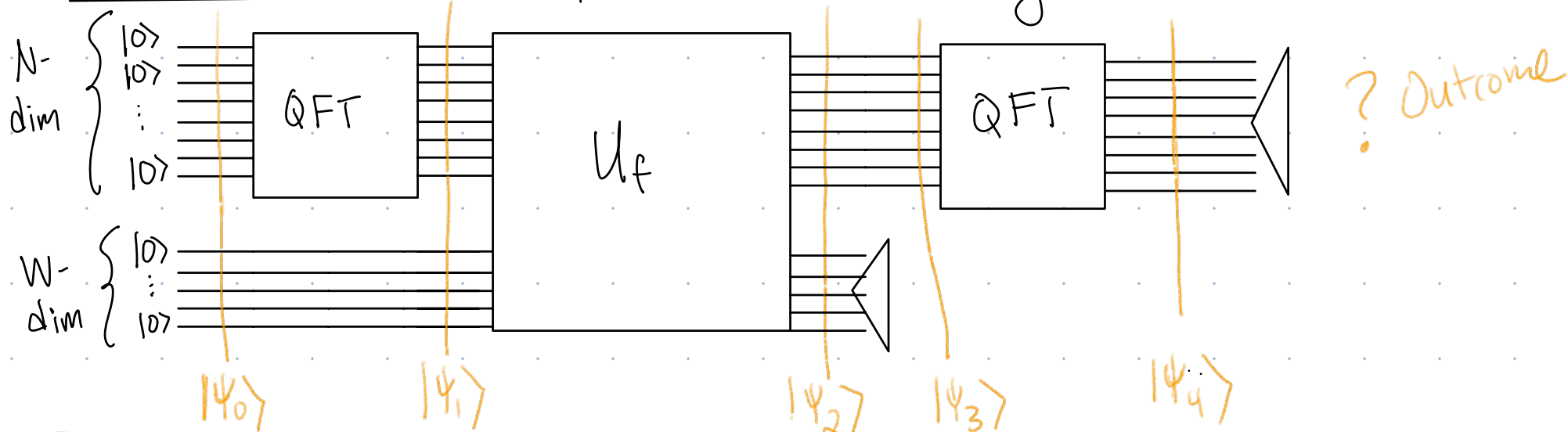
A) $\lceil \log_2 3 \rceil$

B) 3

C) $\log_2 N$

D) N

Quantum Circuit for Period Finding

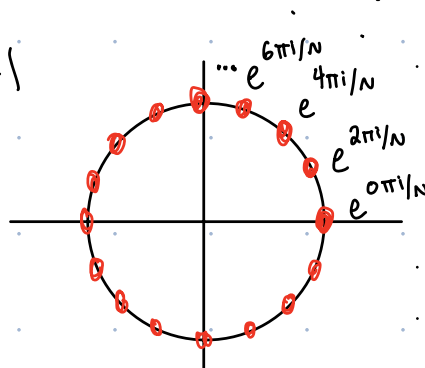


QFT: Quantum Fourier Transform

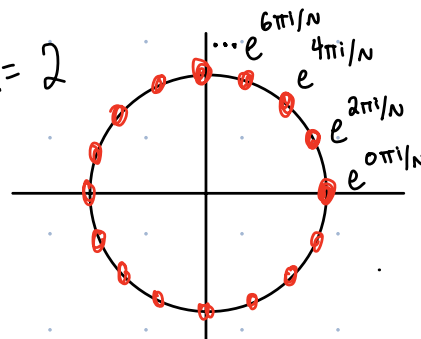
U_f :

Phase of
each s.b.
state in
superposition

$X=1$



$X=2$



$QFT|0\rangle = ?$

go/fourier-transform

Reminder: $U \left(\sum_i a_i |i\rangle \right) = \sum_i a_i U |i\rangle$

$\triangleleft$ = Standard
basis
measurement

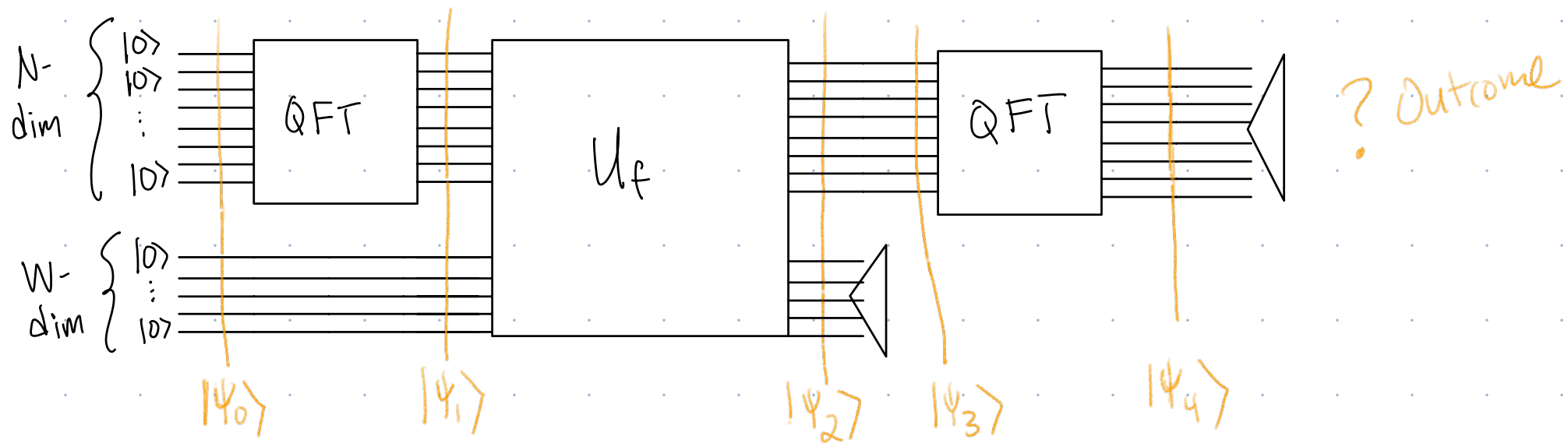
- Measure entire state:

$$|\psi\rangle = \sum_{x=0}^{N-1} a_x |x\rangle \Rightarrow \text{Prob of outcome } |x^*\rangle \text{ is } |a_{x^*}|^2$$

- Partial measurement (B system)

- Factor standard basis states of measured register
- Collapse + renormalize

Group Exercise: Analyze!



$$|\psi_0\rangle =$$

$$|\psi_1\rangle =$$

=

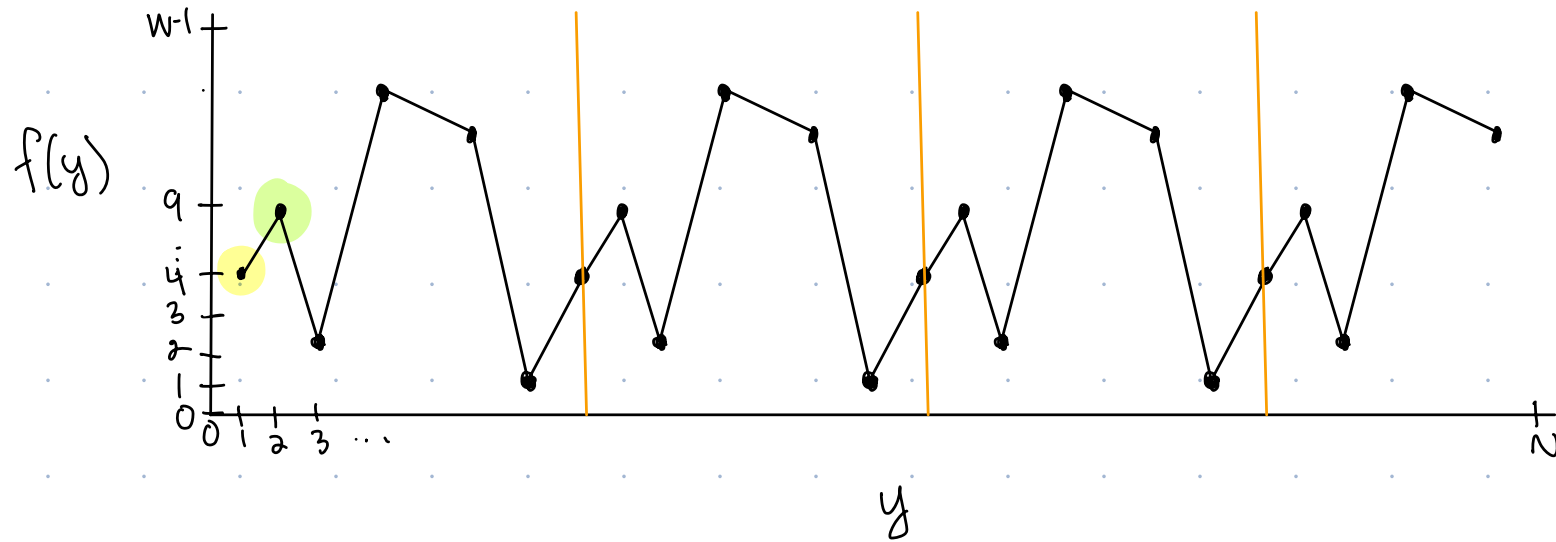
$$|\psi_2\rangle =$$

$$|\psi_3\rangle = \boxed{} \sum_{m=0}^{\boxed{}} |b^* + \boxed{}\rangle$$

(outcome of partial measurement is $f(b^*)$)

For the following example, what does state collapse to if outcome is $|4\rangle$? If outcome is $|9\rangle$? What is pattern if get outcome $|f(b^*)\rangle$?

$r=5$



Suppose measure outcome $|f(b^*)\rangle$:

$$|\psi_4\rangle =$$

=

Plugging In:

$$|\psi_4\rangle =$$

$$|\psi_5\rangle =$$

"

"

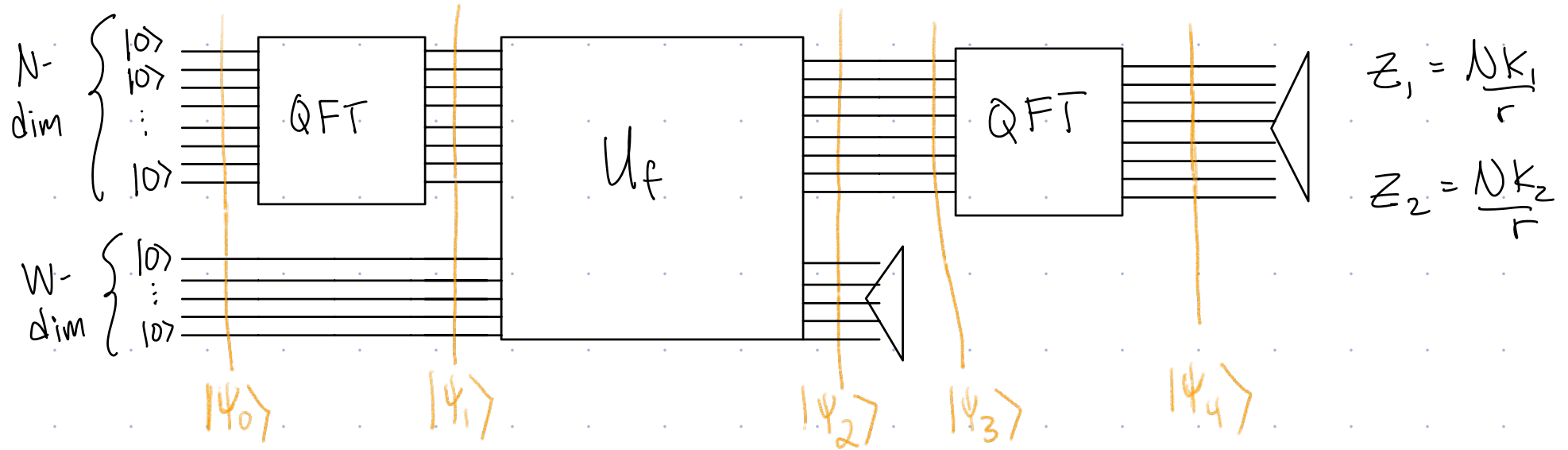
Probability of Outcome $|z\rangle \rightarrow |\text{amplitude of } |z^*\rangle|^2$

Geometric Series Formula:

$$\sum_{m=0}^{t-1} a^m = \begin{cases} t & \text{if } a=1 \\ \frac{1-a^t}{1-a} & \text{if } a \neq 1 \end{cases}$$

Period Finding Algorithm (Shor's Algorithm)

① Run 2 times:



②

③

Query Complexity of Period Finding

Classical: $O(\sqrt{r})$

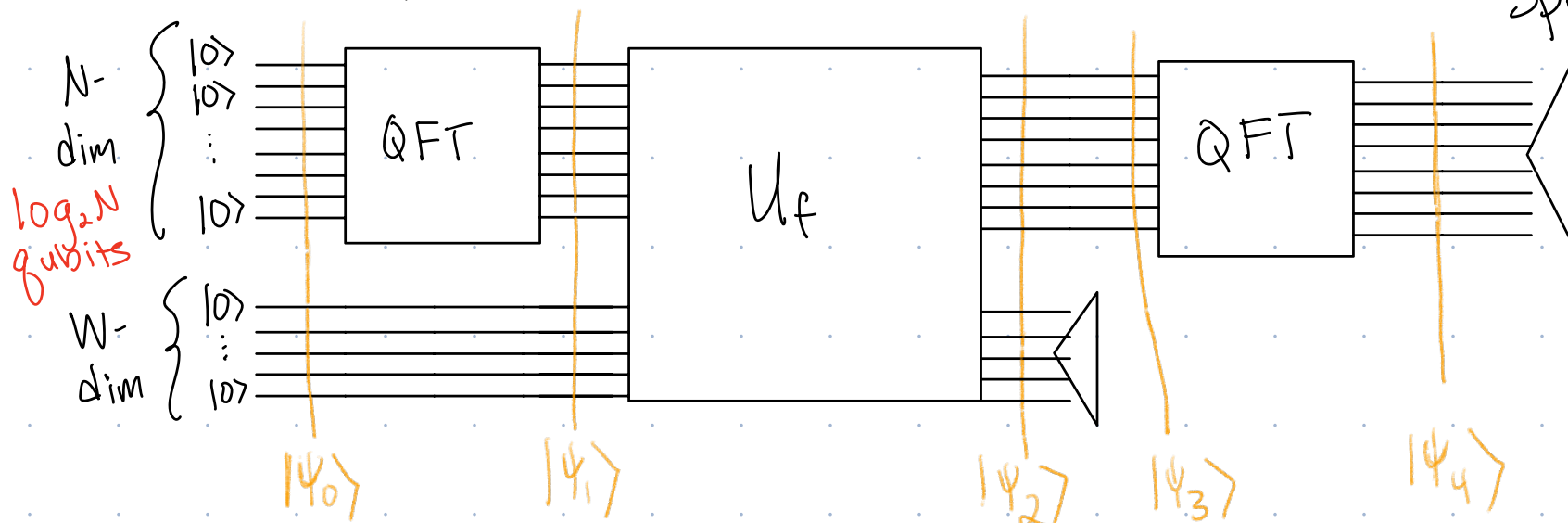
Quantum: $O(1)$

Time Complexity of Factoring

Quantum: $O(\log^2 N)$

Circuit used to factor N .

Nearly Exponential
Speedup!



$O(\log^2 N)$ $O(\log^2 N)$ $O(\log^2 N)$ $\rightarrow O(\log^2 N)$

Classical: $e^{O(\log^{1/3} N)}$

Number Field Sieve

Pesky Detail:

We looked at prob of: $z^* = \frac{kN}{r}$

But if $\frac{N}{r} \notin \mathbb{N}$, z^* is a fraction... see pset.