QuickSort Input: Array A of unique integers Output: Sorted A • If |A|=1: Return A · PivInde randomly chosen index, with value pivVal & Preprocessing · Partition (A, pivInd) & Dividing · QuickSort (AL) (Conquering GuickSort (Ap)) After Partition; Unsorted PivVal unsorted ALALAR Value correct value > pivVa < pivVal sorted location Key Pts · Partition takes O(IAI) time • If pivVal is Z; (ith smallest element of A) after Partition, pivVal is at position i.



Lucky vs. Unlucky Pivot Choices 1. Suppose you get <u>lucky</u> at every recursive call or QuickSort. $T(n) = \begin{cases} O(i) & \text{if } n \leq 1 \\ T(\frac{n}{2}) + T(\frac{n}{2}) + O(n) = aT(\frac{n}{2}) + O(n) \end{cases} f O(n \log n)$ Great! 2. Suppose you get unlucky at every recursive call of QuickSort. $\int (n_{x}) = O(n^{2})$ $T(n) = \int O(1) \quad n \in I$ $T(n) = \int O(1) \quad n \in I$ $T(n-1) + O(n) \quad J = D Expand + Hope$ Awful

٠	Partition (A, pivInd, pivVal)		
٠	·Swap pivot with A[1]		
٠	· Current « 2		
•	· While current = A1:	• •	
٠	XIE Accurrent 7×pivVal: Every element of array	• •	• •
	[Swap A[current], pivVal is compared to pivot.		
	Swap A[pivInd+1], pivVal		
	Current ++ · · · · · · · · · · · · · · · · · ·	• •	• •
•	Current ++ Strategy: count the # of times & is run		
	Strategy: count the # of times & is run over the whole	· ·	· ·
•	Current ++ Strategy: count the # of times * is run over the whole algorithm	· · ·	· · ·
•	Current ++ Strategy: count the # of times & is run over the whole algorithm unsorted pivVal unsorted	· · ·	· · ·
•	Current ++ Strategy: count the # of times * is run over the whole algorithm unsorted pivVal unsorted AL AR	· · · · · · · · · · · · · · · · · · ·	
· · ·	After Partition: Unsorted pivVal unsorted AL AR Value 7 pivVal value 7 pivVal pivVal	· · · · · · · · · · · · · · · · · · ·	

Analyzing Average Kuntime 1. Determine Sample Space S= set of all possible Sequences of random events that might occur over the course of alg. ex: QuickSort -> S=set of possible sequences of pivot choices that the alg might make What is the sample space if QuickSort is run ОИ 857 A) S= 38, 5, 73 B) S = All possible permutations of §8,5,79 C) S = Power set of {8,5,73 (set of all subsets of 28,5,7}) (D) = 7(7), (8, 5), (8, 7), (5, 8), (5, 7)

Z 5 Ŧ 1/3 8. 13 r 8vs 5, 8us7 (\mathbf{c}) 7 vs 5 E 7 15 81 .5 8 5 58 7 7 8 8 5 5 7 Ł 5 12 711/2 7 15 S 6 5 1/6 $S = \{7, (8,7), (8,5), (5,8), (5,7)\}$ R(8,7) = 3R(7) = 2

Analyzing Average Runtime 2. Create « Kandom Variable that maps each element of the sample space to $R: S \rightarrow \mathbb{R}$ a number of times Kg ex: QuickSort: R(U) = [# of comparisons] of 2 elements of our array if pivot sequence J is chosen. 3. Take <u>Expectation value</u> of R to get average run time prob. of J ocquing $\mathbb{E}[\mathbb{R}] = \mathbb{Z} \quad \mathbb{P}(\sigma) \cdot \mathbb{R}(\sigma)$ $e_{X}: [8|5|7] (7) (57) (58) (85)$ (87) $+1 \cdot 3 = 2\frac{2}{3}$ $= \frac{1}{2} \cdot 2 + \frac{1}{6} \cdot 3$ + - - 3 + - 3