

## Learning Goals

- Describe divide + conquer structure ✓
- Describe methods for D & C for
  - proving correctness
  - analyzing time complexity (runtime) ←

## Exit Tickets

# Divide & Conquer

## Merge Sort

**Input** : Integer array  $A$  of length  $n$

**Output**: Sorted array

// Base Case

1 **if**  $n == 1$  **then**

2 | return  $A$ ;

3 **end**

// Divide and Conquer

4  $A_1 = \text{MergeSort}(A[1 : n/2])$ ;

5  $A_2 = \text{MergeSort}(A[n/2 + 1 : n])$ ;

// Combine

6  $p_1 = p_2 = 1$ ;

7 **for**  $i=1$  **to**  $n$  **do**

8 | **if**  $A_1[p_1] < A_2[p_2]$  **then**

9 |  $A[i] = A_1[p_1]$ ;

10 |  $p_1++$ ;

11 **else**

12 |  $A[i] = A_2[p_2]$ ;

13 |  $p_2++$ ;

14 **end**

15 **end**

# Divide & Conquer Structure

## Merge Sort

**Input** : Integer array  $A$  of length  $n$

**Output**: Sorted array

// Base Case

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11 else

12 |  $A[i] = A_2[p_2]$ ;

13 |  $p_2++$ ;

14 end

15 end

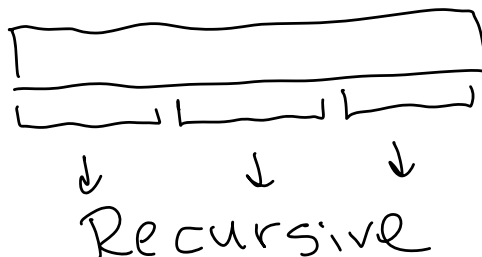
← Base case

← Preprocessing (none in MergeSort)

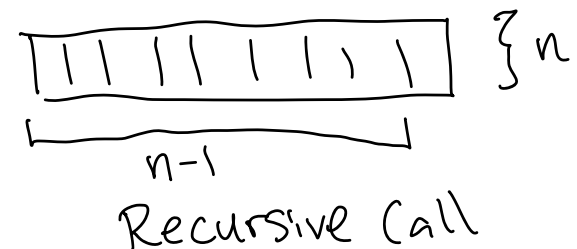
← Divide input into equal sized pieces + recursively solve each piece

← Combine solutions from each recursive call to get ultimate solution

Is D&C



Not D&C



# MergeSort

**Input** : Integer array  $A$  of length  $n$

**Output**: Sorted array

// Base Case

```
1 if  $n == 1$  then  
2   return  $A$ ;  
3 end  
  // Divide and Conquer  
4  $A_1 = \text{MergeSort}(A[1 : n/2])$ ;  
5  $A_2 = \text{MergeSort}(A[n/2 + 1 : n])$ ;  
  // Combine  
6  $p_1 = p_2 = 1$ ;  
7 for  $i=1$  to  $n$  do  
8   if  $A_1[p_1] < A_2[p_2]$  then  
9      $A[i] = A_1[p_1]$ ;  
10     $p_1++$ ;  
11  else  
12     $A[i] = A_2[p_2]$ ;  
13     $p_2++$ ;  
14  end  
15 end
```

How to prove correct?

Strong Induction → See review videos  
(other options?)

How to analyze runtime?

Recurrence Relation

$T(n)$  = runtime on input of size  $n$

$$T(n) = \begin{cases} O(1) & \text{if } n \leq 1 \\ 2T(n/2) + O(n) & \text{else} \end{cases}$$

↓ Solve

$$T(n) = O(n \log n)$$