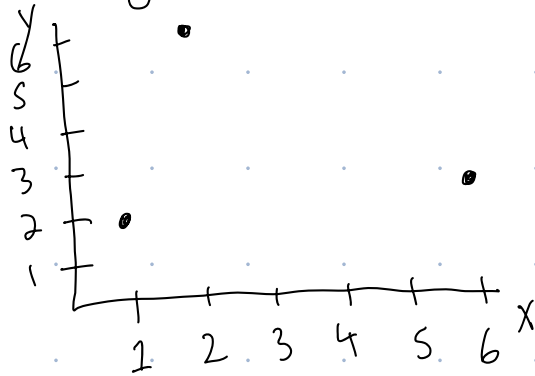


Closest Points Problem

Input: Array of 2-D points:

$P = [(1, 2) \mid (2, 6) \mid (6, 3) \mid \dots]$



Output: Distance b/t 2 closest points

$$\hookrightarrow d(p_i, p_j) = \sqrt{(x_i - x_j)^2 + (y_i - y_j)^2}$$

Applications:

- air traffic control
- robotics
- stereo imaging

Algorithms + Ethics

Algorithm is essentially a mathematical object.

But once it gets implemented for a particular task, has ethical implications.

Assume:

x, y coordinates are
unique for each pt

~~$[(1, 2) \mid (1, 6)]$~~

$[(1, 2) \mid (2, 1)]$

Ethical Matrix (O'Neil + Gunn)

Where travel + why?

Air Traffic
Control
Improvement

Harm?
Benefit?



Choice to use?
Are users informed
enough to understand
meaningfully take
responsibility for use?



Unfair treatment
of different
groups?



Stakeholders	Well-Being	Autonomy	Justice
Airplane Passengers	Safety, Cheaper (?) Environment	A-transport Zoom	Equal access (rich/poor) (geographic)
Airline Cos	\$, safer	No choice (international)	Difference by country, by region
Employees, ATC	Automation → No job Efficiency, less stress	No choice safer	Who is to blame if error
Civilians, Bystanders	Noise pollution, Make war easier, environment, help economy/tourists	No choice, opt out (vote)	(un) Equal impact on environment, near airport more affected

Army			
People of color			
Non-binary folks			

Ethical matrix does not tell you what to do!

Tool for thinking about consequences, both + and -.

→ How can I mitigate negatives? → Modify → Aid

→

Closest Points 2D

- Before designing a sophisticated algorithm, try to benchmark
- Want better than "Brute Force"
 - Can't do better than 1D

Brute Force (check every pair)

- $\text{min} \leftarrow \infty$ $O(1)$
- for $i \leftarrow 1$ to $n-1$ $O(n)$
 - for $j \leftarrow i+1$ to n $O(n)$ $\} O(n^2)$
 - if $\text{dist}(p_i, p_j) < \text{min}$, then $\text{min} \leftarrow \text{dist}(p_i, p_j)$ $O(1)$
- return min

Closest Pts 1D

- $\text{Sort}(P)$ $O(n \log n)$
- $\text{min} \leftarrow \infty$ $O(1)$
- for $i = 1$ to $n-1$: $O(n)$
 - if $\text{dist}(p_i, p_{i+1}) < \text{min}$, then $\text{min} \leftarrow \text{dist}(p_i, p_{i+1})$
- return min

5	9	0	20	17
---	---	---	----	----

0	5	9
---	---	---

0	5	9	17	20
---	---	---	----	----

17	20
----	----

- $\text{Sort}(P)$ $O(n \log n)$
- $O(1) \leftarrow \} O(n) \} O(n \log n)$
 $O(1)$

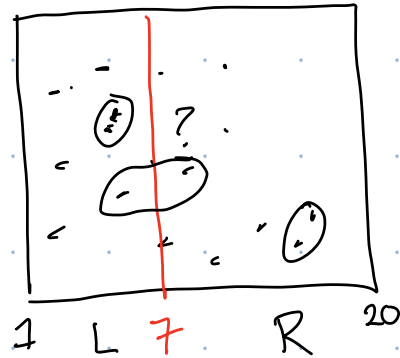
Ethical Matrix: Race? Include the right stakeholders/how many?
Weakness? When to use? What to do after fill out?

CloPts(P) (Divide + Conquer 2D Closest Pts)

Base Case: Later

Sort by x

Divide:



(1,5)	(2,2)	(4,0)	(10,20)	(12,3)	(20,6)
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L

R

midline: 7

Conquer:

$S_1 \leftarrow \text{CloPts}(L)$

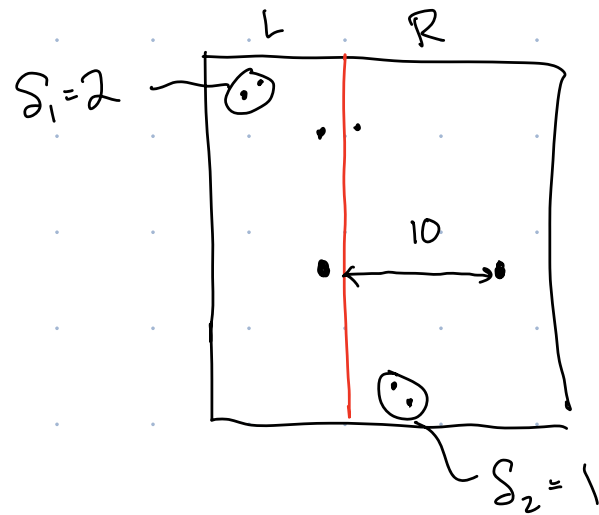
$S_2 \leftarrow \text{CloPts}(R)$

$S \leftarrow \min(S_1, S_2)$

Combine

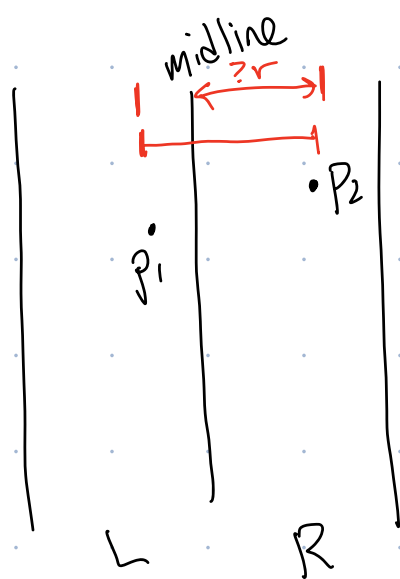
Combine:

Let's think about this:



Lemma: If p_2 in R has difference in x-coord. of more than $r = \delta$ from midline, then for any point $p_1 \in L$, $d(p_1, p_2) > r = \delta$

Proof Sketch: (Basic ideas, but need more English in real proof)



$$d(p_1, p_2) = \sqrt{(x_1 - x_2)^2 + (y_1 - y_2)^2}$$

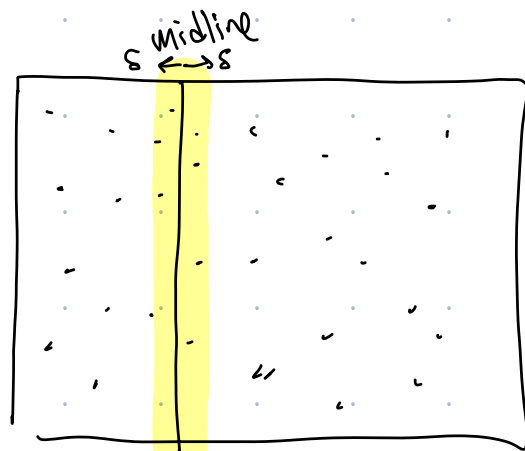
$$(x_1 - x_2)^2 = \left(\underbrace{\text{diff in x-coord from } x_2 \text{ to midline}}_{\text{diff "}} + \underbrace{\text{" midline to } x_1}_{> 0} \right)^2 > r^2$$

$$(y_1 - y_2)^2 \geq 0$$



$$\sqrt{(x_1 - x_2)^2 + (y_1 - y_2)^2} \geq \sqrt{r^2}$$

$$d(p_1, p_2) \geq r$$



We want to check points in this region,
Looks like a line!

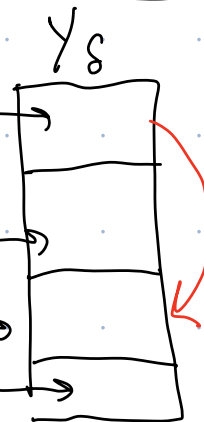
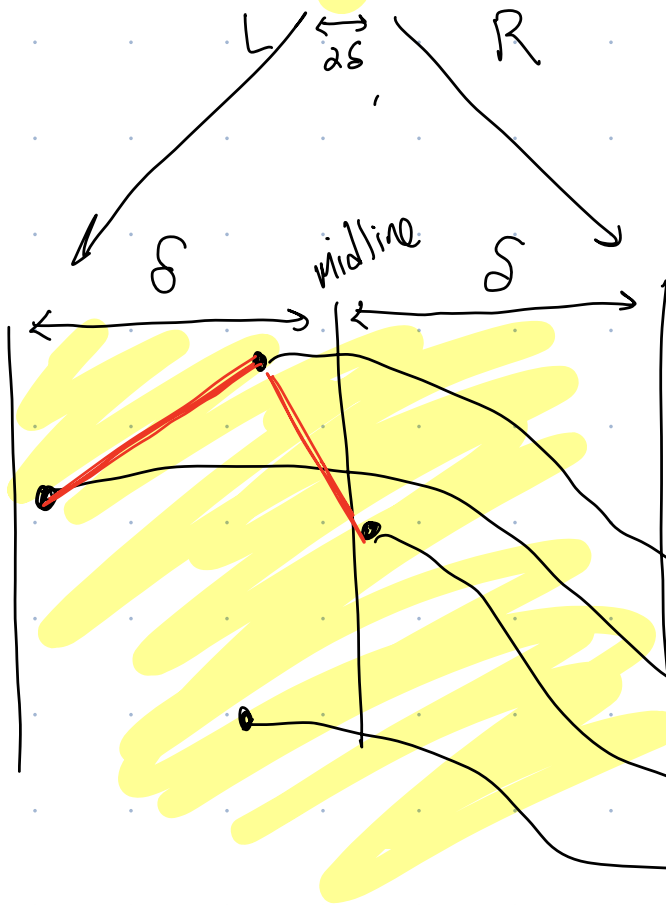
Combine Step

$Y_\delta \leftarrow$ y-sorted list of pts in P
within δ of midline

For $p \in Y_\delta$

- check distance from p to next 2+ pts
- Save if smallest found

Return smallest distance found

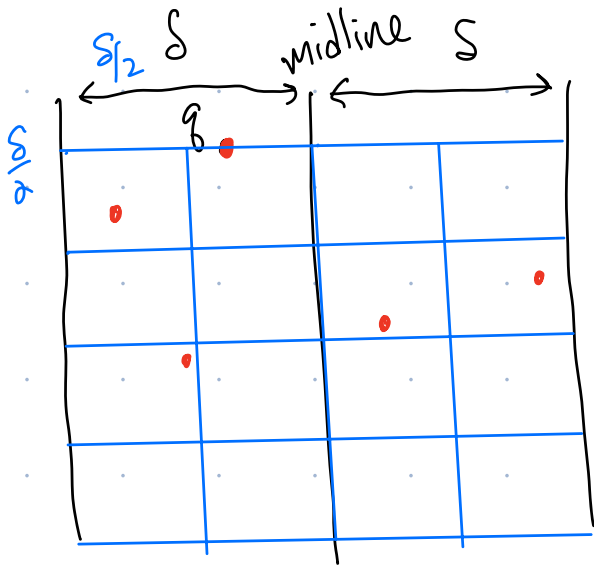


What goes here?

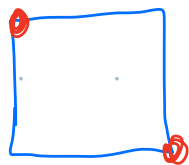
Lemma: Only need to look at next 7 pts in Y_S

Pf: Imagine dividing region w/in S of midline into $\frac{S}{2} \times \frac{S}{2}$ squares starting at current pt (q) . Each square can contain at most one point. To see this, for

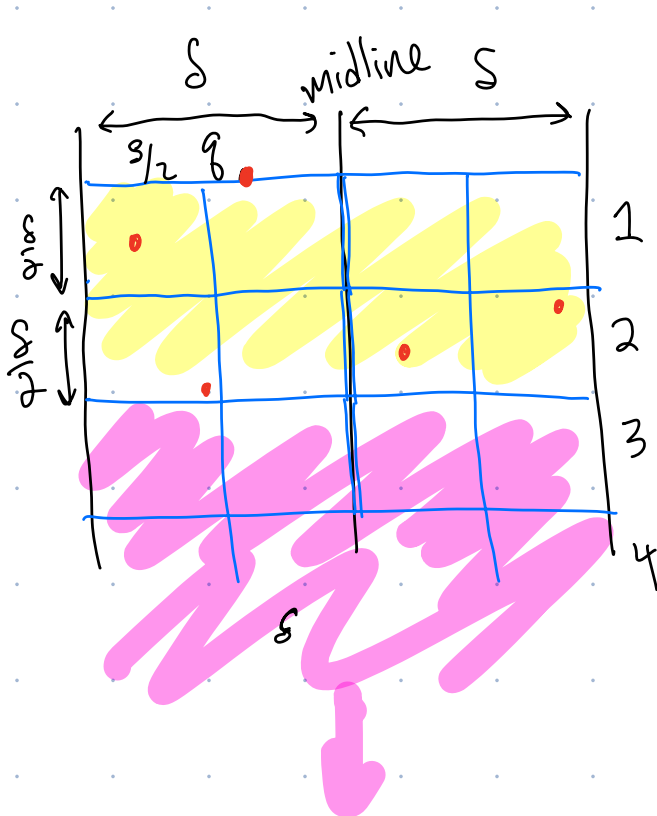
contradiction, suppose there are 2 pts in a square. The pts have largest distance when on opp. corners. In that case, their distance $> S$



$$\sqrt{\left(\frac{S}{2}\right)^2 + \left(\frac{S}{2}\right)^2} = \sqrt{\frac{S^2}{4} + \frac{S^2}{4}} = \sqrt{\frac{S^2}{2}} = \frac{S}{\sqrt{2}} < S$$



Each box is in L or R, so any 2 pts in a box must have distance $\geq S$, a contradiction



Pts in rows $3+$ have distance at least δ from q , b/c diff in y -coordinate is at least δ .

So only 1st 2 rows are relevant

There are at most 7 pts other than q in 1st 2 rows, b/c only 8 box.

