

Goals

- Write Proof By Contradiction
- Write Strong inductive proof.

Test Info

- All topics up to iff proofs on 3/6
- All PSets up to 4
- Monday : review . Topics → Canvas Discussion until Saturday
- Wed @ ~~noon~~^{10am} → Frid @ 6 pm
- Reserve 3 hours
- 1 page, single sided handwritten notes "cheat sheet"

Quiz on Canvas

Short PSet 5

Goals

- Describe & write proof by contradiction
- Describe & write proof by strong induction

Proof by Contradiction

Use: any statement P

Proof needs to do two things

①	(Direct)	$\Gamma P \rightarrow Q$	}	Most common
②	(Direct)	$\Gamma P \rightarrow \Gamma Q$		
$\therefore P$				

Structure

For contradiction assume $\neg P$

Explain explain... Q

Explain explain... $\neg Q$, a contradiction.

Thus, P is true

When start, don't know what Q is... you need to keep your eye out for what might be the contradiction.

Q: Prove $\sqrt{2}$ is irrational

* Not of form $P \rightarrow Q$

A: For contradiction, suppose $\sqrt{2}$ is rational. Then
 $\exists a, b \in \mathbb{Z} : \frac{a}{b} = \sqrt{2}$ where the fraction
 is fully simplified, so $\nexists c \in \mathbb{Z} : c|a \wedge c|b$.
 Squaring both sides, we have

$$a^2 = 2b^2.$$

Thus $2|a^2$. But we've previously proved this
 implies $2|a$. This means $\exists m \in \mathbb{Z} : 2m = a$. Plugging
 in, we have

$$4a^2 = 2b^2.$$

Dividing by 2, we get

$$2a^2 = b^2.$$

But this means $2|b^2$, and so $2|b$. But
 this means $2|a$ and $2|b$, which contradicts the fact
 that $\frac{a}{b}$ is fully simplified. \square

Q: Prove: $\nexists x, y \in \mathbb{Z} : x^2 = 4y^2 + 2$

Try: $\neg \exists x, y \in \mathbb{Z}: x^2 = 4y + 2$ (1st step: is x^2 even or odd)
 for contradiction, there

Suppose \wedge exist $x, y \in \mathbb{Z}: x^2 = 4y + 2$. Then x^2 is even, so
 x is even. Thus $\exists m \in \mathbb{Z}: x = 2m$. Plugging in, and
 solving for y , we have

$$y = \frac{4m^2 - 2}{4} = m^2 - \frac{1}{2}$$

Since m is an integer, $m^2 - \frac{1}{2}$ is not an integer, a
 contradiction.

Other ways to get contradictions:

To Prove $P \implies$ Prove $\neg P \rightarrow P$

To Prove $P \rightarrow Q \implies$ Prove $(P \wedge \neg Q) \rightarrow Q$

$$\neg(P \rightarrow Q) \equiv P \wedge \neg Q$$

To Prove $P \implies \neg P \rightarrow$ something obviously false

Strong Induction Proof Structure

Set-up

Let $P(n)$ be the predicate _____ We will prove
 $P(n)$ is true for all $n \geq$ b.c.
 \uparrow
base case

Base-case

Base-case: We prove $P(\text{b.c.})$ is true. ... (Sometimes multiple b.c.)
 $P(0), P(1)$

Inductive case

Inductive case: Assume for strong induction that $P(j)$ is true
for all j such that $\text{b.c.} \leq j \leq k$.

Therefore $P(k+1)$ is true

Conclusion

By strong induction, we conclude $P(n)$ is true for
all $n \geq$ b.c.

Metaphor:

