

Announcements

- I am away Wed./Fri
 - Friday: in-class quiz (Prof Andrews)
 - I will post video lectures to website.
 - You may want to come to class and watch discuss together
- Self-Grade/Reflections: Due by 3 pm Wed to my mailbox or submitted on Canvas
 - ↑ outside MBH 633
- Extra Office Hours on Wed, Thurs, Friday
 - Book at [Calendly.com/skimmel](https://calendly.com/skimmel)
 - Will use ZOOM. I will send Meeting ID
 - Go to [Middlebury.zoom.us](https://middlebury.zoom.us)
 - Click Join & put in ID
- Extra office hours M, T drop-in

$\forall x \exists y$
 ↑ ↑
 Nemesis gets to choose from any in domain to try to make false
 You get to choose one specific y based on x from nemesis and prove true

$P(x, y) = x + y = 10$
 Can get rid of y in predicate using quantifier:
 $P(x) = \forall y, x + y = 10$

Direct Proof

- Usually used to prove implication like $P \rightarrow Q$.

Structure:

Assume P . Explain, explain, ... explain. Therefore Q
 "By definition..."

- Also used to prove universal implication: $\forall x (P(x) \rightarrow Q(x))$
 \uparrow
 x in some Domain

Structure:

Let x be [any arbitrary] element of the domain

Assume $P(x)$

Explain, explain, explain.

Therefore $Q(x)$

Q: Use a direct proof to show: For all $a, b, c \in \mathbb{Z}$, if $a|b$ and $b|c$ then $a|c$.

($a|b$ means a divides b , that $\exists e \in \mathbb{Z}: ae = b$.)

Let $a, b, c \in \mathbb{Z}$. Assume $a|b$ and $b|c$. This means $\exists e, f \in \mathbb{Z}$ such that $b = ae$ and $c = bf$. Then $c = bf = (ae)f = a(ef)$.

But this means $a|c$, since $c = ak$, for $k = ef$, an integer.

$P \iff Q$ If and only if

$(P \rightarrow Q) \wedge (Q \rightarrow P)$ logically equivalent to $P \iff Q$

\Downarrow

Structure

For the forward direction, [Proof of $P \rightarrow Q$]

For the backward direction [Proof of $Q \rightarrow P$]

Direct Proof - Proof by Contrapositive

- Usually used to prove implication like $P \rightarrow Q$.

Recall: $P \rightarrow Q$ is logically equivalent to $\neg Q \rightarrow \neg P$

Structure:

We prove the contrapositive

Assume $\neg Q$. Explain, explain, ... explain. Therefore $\neg P$

- Also used to prove universal implication: $\forall x (P(x) \rightarrow Q(x))$
 \uparrow
 x in some Domain

Structure:

Let x be [any arbitrary] element of the domain

We prove the contrapositive.

Assume $\neg Q(x)$

Explain, explain, explain.

Therefore $\neg P(x)$

(See example on hw!)

What if statement is not of the form $P \rightarrow Q$?

What if just have P ?

Suppose you can show

Proof needs to do two things

①	(Direct)	$\Gamma P \rightarrow Q$
②	(Direct)	$\Gamma P \rightarrow \neg Q$
		$\therefore P$

P	Q	ΓP	ΓQ	$\Gamma P \rightarrow Q$	$\Gamma P \rightarrow \neg Q$
T	T	F	F	T	T
T	F	F	T	T	T
F	T	T	F	T	F
F	F	T	T	F	T

Structure: (Prove P)

" For contradiction, assume ΓP "

or
 " We proceed by contradiction. We assume $\neg P$. "

⋮

Therefore, Q

However

⋮

Therefore, $\neg Q$, a contradiction. Thus, P .