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Ways to Represent Graphs in Computer
Adjacency Matrix

|  | 1 | 2 | 3 | 4 |
| :--- | :--- | :--- | :--- | :--- |
| 1 | 0 | 0 | 1 | 1 |
| 2 | 0 | 0 | 0 | 1 |
| 3 | 1 | 0 | 0 | 1 |
| 4 | 1 | 1 | 1 | 0 |



Store as array $A$ in memory. Can learn $A[i, j]$ in $O(1)$ time.

Which adjacency matrix represents this graph?

| 0 | 1 | 0 | 0 |
| :--- | :--- | :--- | :--- |
| 1 | 0 | 1 | 1 |
| 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | 0 |

A


| 0 | 1 | 1 | 1 |
| :--- | :--- | :--- | :--- |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 |
| 1 | 0 | 0 | 0 |

B

| 0 | 1 | 0 | 0 |
| :--- | :--- | :--- | :--- |
| 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 1 |
| 0 | 0 | 1 | 0 |

C

Adjacency List

| Vertex | Adjacent Vertices |
| :---: | :---: |
| 1 | 3,4 |
| 2 | 4 |
| 3 | 1,4 |
| 4 | $1,2,3$ |



Store as an array of lists
can access $A[3]$, (the list) in $O(1)$ time, but then to go through list takes time $O(L)$ where $L$ is length of list. Can learn $A[3]$. length io $O(1)$ time.
Edge List
Can represent graph as list of edges, but worst case time complexity bad for most applications
def: The degree of a vertex is the number of adjacent edges.
def: A vertex $v_{1}$ is adjacent to vertex $v_{2}$ if connected by an edge

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How would you represent a

- directed graph?
- graph with self-loops?
- graph with multiedges?
- graph with weighted edges?


Using Adjacency Matrix / Adjacency List?
Give representations of this graph using both approaches:

| $v$ | List |
| :--- | :--- |
| 1 | $(2,1 / 3),(3,1 / 3),(4,1 / 3)$ |
| 2 | $(1,1 / 2),(4,1 / 2)$ |
| 3 | $(4,1)$ |
| 4 | $(4,1)$ |

