1. Create a recurrence relation for the worst case runtime of the following algorithm for binary search when \( f - s + 1 = n \). You may assume \( n \) is a power of 2. Use the iterative method to solve the recurrence relation.

**Algorithm 1: BinarySearch\((A, x, s, f)\)**

**Input** : Sorted (in increasing order) array of integers \( A \), an integer \( x \) that occurs in the array, a starting index \( s \) and an ending vertex \( f \)

**Output**: An index \( i \) such that \( A[i] = x \).

1. if \( s == f \) then
2. return \( s \);
3. end
4. \( \text{mid} = \lfloor (s + f)/2 \rfloor \);
5. if \( A[\text{mid}] < x \) then
6. return \( \text{BinarySearch}(A, x, \text{mid} + 1, f) \)
7. else
8. return \( \text{BinarySearch}(A, x, s, \text{mid}) \)
9. end

**Solution**  Let \( T(n) \) be the run time of the algorithm, when \( f - s + 1 = n \). Then

\[
T(n) = T(n/2) + C
\]  (1)

where \( C \) is some constant. The initial conditions are

\[
T(1) = B
\]  (2)

where \( B \) is some constant.

Then using the iterative approach, we have

\[
T(n) = T(n/4) + C + C = T(n/4) + 2C = T(n/8) + C + 2C = T(n/8) + 3C.
\]  (3)

We see the pattern is

\[
T(n) = T(n/2^k) + k \times C.
\]  (4)

The base case is when \( 2^k = n \), which happens when \( k = \log_2(n) \). Plugging this in, we have

\[
T(n) = T(1) + \log_2(n) \times C = O(\log_2(n)).
\]  (5)
2. Let $K(n)$ be the size of the set of $n$-digit numbers that have an even number of 0’s. Create a recurrence relation for $K(n)$. What is $K(3)$? (Hint 0: remember zero is even. Hint 1: think about the possible options for the value of the final digit. of the number. Hint 2: The size of the set of numbers that don’t have an even number of 0’s is the total number of elements minus the set of numbers that do have an even number of 0’s.)

Solution The base case is $K(1) = 9$ since the only way you can have a 1-digit number with an even number of 0s is to have no 0s.

For the recurrence relation, We can break up the set of $n$-digit numbers with an even number of 0s into those that have a 0 in the final position, and those that don’t. For those that have a final 0, there must be an odd number of 0s in the $n-1$ digits preceding the final digit. The number of $(n-1)$-digit numbers with an odd number of 0s is the total number of $(n-1)$ digit numbers (which is $10^{n-1}$), minus the number of $(n-1)$-digit numbers with an even number of 0s (which is $K(n-1)$). Thus the number of $n$-digit numbers with an even number of 0s and the last digit equal to 0 is

$$10^{n-1} - K(n-1).$$

(6)

In the case that the final digit is not a zero, then for there to be an even number of 0s, we must have an even number of zeros in the first $n-1$ digits. There are $K(n-1)$ such numbers. Since we have a choice of 9 digits for the last digit, there are $9 \times K(n-1)$ numbers that don’t have a final digit 0. Thus the number of $n$-digit numbers with an even number of 0s and the last digit equal to 0 is

$$9K(n-1).$$

(7)

Putting it all together, we have

$$K(n) = 9K(n-1) + 10^{n-1} - K(n-1) = 8K(n-1) + 10^{n-1}.$$  

(8)

Thus

$$K(3) = 8K(2) + 10^2 = 8(8K(1) + 10^1) + 10^2 = 8 \times 8 \times 9 + 110 = 686.$$  

(9)

3. Create a recurrence relation for the number of ways a person can climb $n$ stairs if the person can take one stair or two stair at a time. How many ways can this person climb a flight of 8 stairs?

Solution Let $S(n)$ be the number of ways the person can take $n$ stairs. For the initial conditions, we have

$$T(1) = 1$$

(10)

because there is only one way to take one stair, and

$$T(2) = 2$$

(11)
because the person could either take the 2 stairs at once, or could go up one at a time.

Now for the recurrence relation, if the person takes \( n \) stairs, they could either take the last stair on its own, or they could take the last two stairs together.

If they take the final stair as one step, there are \( T(n - 1) \) ways that they could take the first \( n - 1 \) stairs to get there, so there are \( T(n - 1) \) ways of taking the final stair alone.

If they take the final two stairs together, there are \( T(n - 2) \) ways that they could have gotten to the \( (n - 2) \)th stair, and then one way that they can take the final two together.

Thus

\[
T(n) = T(n - 1) + T(n - 2). \tag{12}
\]

So

\[
\]
\[
= T(5) + T(4) + 2(T(4) + T(3)) + T(3) + 2
\]
\[
= T(5) + 3(T(4) + T(3)) + 2 = T(4) + T(3) + 4(T(3) + T(2) + T(2) + T(1)) + 2 = \tag{13}
\]

...I got bored, but you see where this is going, hopefully!